## CTDB Status - Clustered Samba Growing Up

### SNIA Storage Developer Conference 2011

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Samba Team / SerNet

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```
[root@node0 ~]# ctdb status
Number of nodes:3
pnn:0 192.168.46.70
                       OK (THIS NODE)
pnn:1 192.168.46.71
                       0K
pnn:2 192.168.46.72
                       0K
Generation:2061920893
Size:3
hash:0 lmaster:0
hash:1 lmaster:1
hash:2 lmaster:2
Recovery mode:NORMAL (0)
Recovery master:1
[root@node0 ~]#
```

## Thank you very much!

## Introduction and History



SerNet





- Prerequisite: cluster file system
- ▶ all-active ⇒ all nodes act as **one** CIFS server
- ► IPC: messaging
- ▶ IPC: sessions, connections, open files, locks, ... (TDB databases)
- ▶ Persistent data: secrets, registry, id-map, ... (TDB databases)
- ► TDB: small, fast, key-value database with record locks and memory mapping

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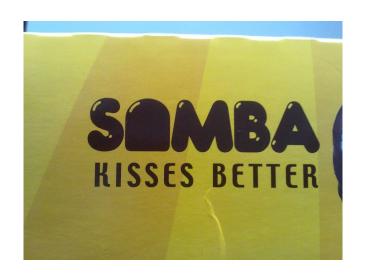
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## ... with CTDB ©



SerNet



- ▶ is a very special clustered database implementation (may lose data...)
- is a samba-inter-node-IPC implementation
- is a simple cluster service management software
- makes Samba on a file system cluster appear as a single CIFS/SMB/SMB2 server
- does not require any client changes to access the Samba cluster



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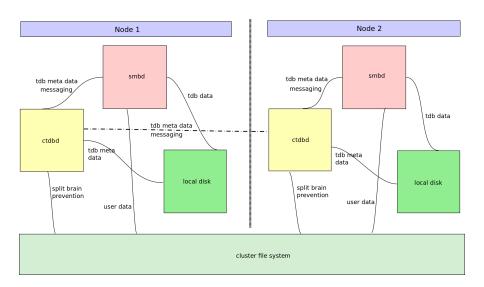
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- ▶ first usable version: 2007 (non-persistent DBs only)
- ▶ today: Ronnie Sahlberg maintainer
- ▶ support for persistent DBs added 2008/2009
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- ▶ code: git://git.samba.org/ctdb.git
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# Recent and Current Improvements

- vacuuming
- persistent transactions
- samba persistent db performance tuning
- tools
- read-only record copies
- further projects





**SerNet** 

- purpose: garbage collection of deleted records
- ▶ was not working well at all:
- under certain workloads, databases grew despite vacuuming
- vacuuming child even crashed under certain conditions
- ▶ ⇒ whole clusters went on strike



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- samba enhancement to schedule upon delete operation
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- several internal polishing tasks



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- lack of global state
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#### Persistent Databases: TODOs

- implement transactions inside CTDB (currently much of the logic is in samba client code)
- correctly handle delete record operations
- implement recoveries for persistent databases differently



# Samba DB Tuning



- ▶ more frequently used persistent DBs too slow in a cluster
- e.g. idmap and registry
- all write operations to persistent database are protected by transactions
- especially expensive in a cluster
- ▶ id-mapping code was rewritten to make creation of ID mappings atomic (2010/2011, Michael Adam)
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# Tools



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- exploit SMB2 durable file handles (Stefan Metzmacher, (Michael Adam))
- exploit new SMB 2.1 and 2.2 features large MTU, multi-credit, leases, re-auth, multi-channel, ... (Stefan Metzmacher, (Michael Adam))

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# Management / Integration



- manages services (samba/winbind/nfs/apache/...): start/stop/monitor
- pluggable extensible event script architecture (/etc/ctdb/events.d/)
- handles IP (re)allocation on public network: fail-over/fail-back
- tickles clients to reconnect in case of fail-over
- When this was created, Linux cluster stack did not have all-active
- But nowadays, pacemaker is getting more popular in distributions
- ▶ All of the above CTDB features are optional



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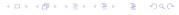


#### Independently of Linux cluster stack

- ► CTDB manages samba
- ► CTDB manages winbindd
- ► CTDB manges public IP addresses

#### As managed resources

- ► CTDB does not manage samba, winbind nor public IPs
- CTDB only provides clustered TDB services
- ► Linux cluster suite (pacemaker) manages CTDB and Samba and Winbind
- ▶ Resource dependency: Cluster FS ⇒ CTDB ⇒ winbindd ⇒ samba



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# Thank you very much!

(Now really... ©)